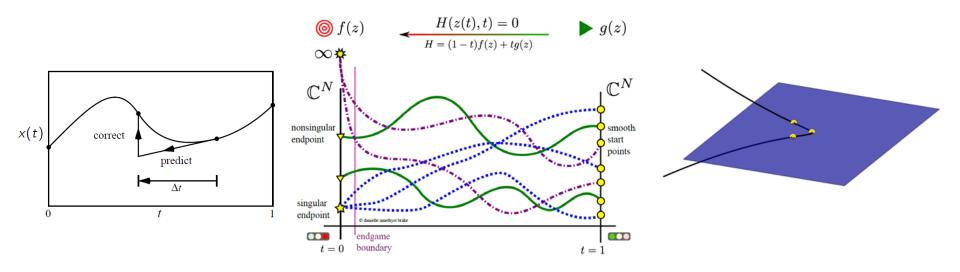
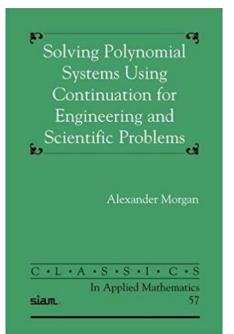
# Foundations of Numerical Algebraic Geometry

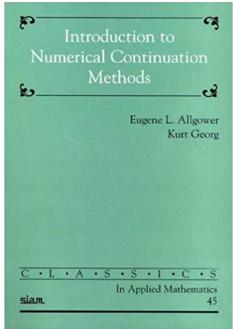


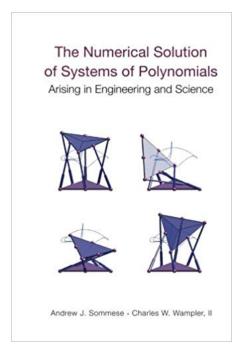
#### Jonathan Hauenstein Nonlinear Algebra Bootcamp September 10, 2018

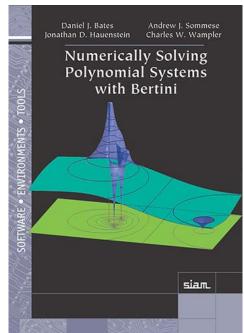
















For a polynomial system  $f: \mathbb{C}^n \to \mathbb{C}^N$ , solve f(x) = 0.





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Maple

```
> solve(x^5 - x + 1);

RootOf(\_Z^5 - \_Z + 1, index = 1), RootOf(\_Z^5 - \_Z + 1, index = 2), RootOf(\_Z^5 - \_Z + 1, index = 3),

RootOf(\_Z^5 - \_Z + 1, index = 4), RootOf(\_Z^5 - \_Z + 1, index = 5)
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RootOf(Z^5 - Z + 1, index = 4), RootOf(Z^5 - Z + 1, index = 5)
```

```
> fsolve(x^5-x+1);
```

-1.167303978





For a polynomial system  $f: \mathbb{C}^n \to \mathbb{C}^N$ , solve f(x) = 0.

#### Maple

```
> solve(x^5 - x + 1);

RootOf(\_Z^5 - \_Z + 1, index = 1), RootOf(\_Z^5 - \_Z + 1, index = 2), RootOf(\_Z^5 - \_Z + 1, index = 3),

RootOf(\_Z^5 - \_Z + 1, index = 4), RootOf(\_Z^5 - \_Z + 1, index = 5)
```

```
> fsolve(x^5 - x + 1); -1.167303978
```

```
> evalf(solve(x^5 - x + 1));
 0.764884433600585 + 0.352471546031726 \text{ I}, -0.181232444469875 + 1.08395410131771 \text{ I}, -1.16730397826142, -0.181232444469875 - 1.08395410131771 I, 0.764884433600585 - 0.352471546031726 I
```





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0.764884433600585 + 0.352471546031726 I, -0.181232444469875 + 1.08395410131771 I, -1.16730397826142, -0.181232444469875 - 1.08395410131771 I, 0.764884433600585 - 0.352471546031726 I
```

#### Bertini

input

```
variable_group x;
function f;
f = x^5 - x + 1;
```

#### finite\_solutions

5

7.648844336005847e-01 -3.524715460317264e-01

7.648844336005849e-01 3.524715460317262e-01

-1.812324444698754e-01 1.083954101317711e+00

-1.167303978261419e+00 -2.220446049250313e-16

-1.812324444698754e-01 -1.083954101317711e+00





For a polynomial system  $f: \mathbb{C}^n \to \mathbb{C}^N$ , solve f(x) = 0.

What does it mean to solve? Some examples include:





For a polynomial system  $f: \mathbb{C}^n \to \mathbb{C}^N$ , solve f(x) = 0.

What does it mean to **solve**? Some examples include:

- Show that a solution exists
  - Computing upper bounds on rank of a tensor

Strassen (1969): rank  $M_2 \le 7$  showing  $\omega \le \log_2 7 < 3$ 

$$\begin{bmatrix} a & b \\ c & d \end{bmatrix} \cdot \begin{bmatrix} A & B \\ C & D \end{bmatrix} = \begin{bmatrix} a \cdot A + b \cdot C & a \cdot B + b \cdot D \\ c \cdot A + d \cdot C & c \cdot B + d \cdot D \end{bmatrix} = \begin{bmatrix} I + IV - V + VII & III + VI \\ II + IV & I - II + III + VI \end{bmatrix}$$

$$egin{array}{llll} I: & (a+d)\cdot (A+D) & & V: & (a+b)\cdot D \\ II: & (c+d)\cdot A & & VI: & (c-a)\cdot (A+B) \\ III: & a\cdot (B-D) & & VII: & (b-d)\cdot (C+D) \\ IV: & d\cdot (C-A) & & & \end{array}$$

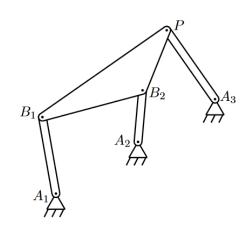




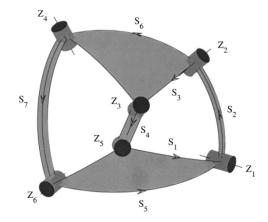
For a polynomial system  $f: \mathbb{C}^n \to \mathbb{C}^N$ , solve f(x) = 0.

What does it mean to **solve**? Some examples include:

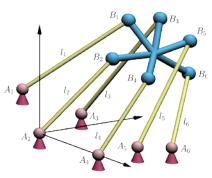
- ▶ Compute all isolated solutions (over  $\mathbb{C}$  or  $\mathbb{R}$ ).
  - Number of assembly configurations



planar pentad SE(2)



spherical pentad SO(3)



Stewart-Gough plaftorm SE(3)

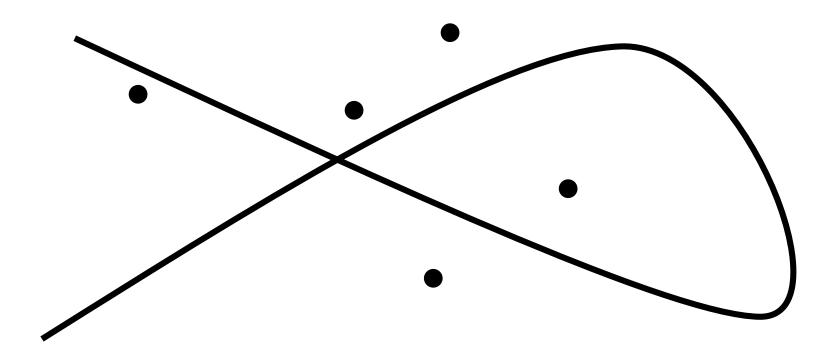




For a polynomial system  $f: \mathbb{C}^n \to \mathbb{C}^N$ , solve f(x) = 0.

What does it mean to **solve**? Some examples include:

Describe all irreducible components.







For a polynomial system  $f: \mathbb{C}^n \to \mathbb{C}^N$ , solve f(x) = 0.

#### Generally speaking:

- Algebraic methods prefer vastly over-determined systems
  - fewer "new" polynomials to compute
  - Bardet-Faugere-Salvy (2004)
- Numerical algebraic geometry prefers well-constrained systems of low degrees with coefficients of roughly unit magnitude
  - ▶ codimension = # equations
  - stable under perturbations





For a polynomial system  $f: \mathbb{C}^n \to \mathbb{C}^N$ , solve f(x) = 0.

What does it mean to numerically solve?





For a polynomial system  $f: \mathbb{C}^n \to \mathbb{C}^N$ , solve f(x) = 0.

What does it mean to numerically solve?

Need two key aspects:

- compute sufficiently accurate numerical approximation
- have an algorithm that can produce approximations of solution to any given accuracy starting from numerical approximation
  - sufficiently accurate depends on the algorithm





What does it mean to numerically solve?

- compute sufficiently accurate numerical approximation
- have an algorithm that can produce approximations of solution to any given accuracy starting from numerical approximation

## Example

$$f(x) = x^2 - 2 = 0$$

 $\triangleright$   $x_0 = 1$  is numerical solution associated with Newton's method





# Overview

$$f(x) = x^2 - 2 = 0$$

 $\triangleright$   $x_0 = 1$  is numerical solution associated with Newton's method

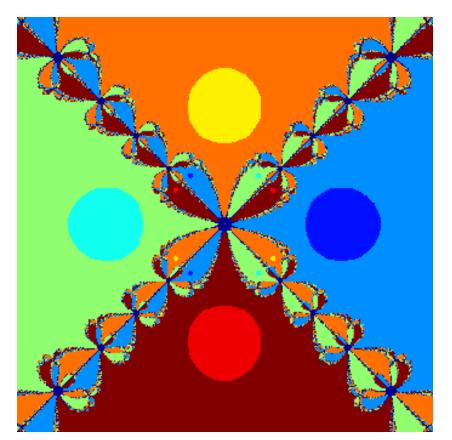
$$x_{k+1} = x_k - Jf(x_k)^{-1}f(x_k)$$





Double-edged sword of Newton's method:

- Qudaratic convergence near nonsingular solutions
- Slow convergence or divergence near singular solutions
- Difficulty away solutions (chaos, limit cycles, etc)



$$f(x) = x^4 - 1$$

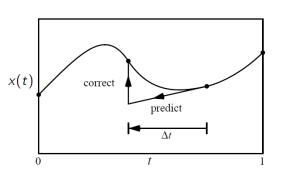


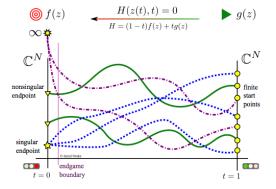


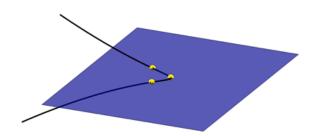
#### Foundations of Numerical Algebraic Geometry:

- Continuation and path tracking
  - Constructing homotopies

- Witness sets
  - Numerical Irreducible Decomposition
  - Other computations using witness sets









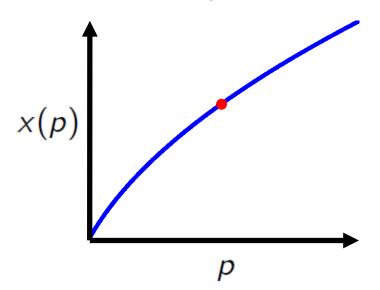


#### Continuation from complex analysis:

- Cauchy (1789-1857), Riemann (1826-1866), Mittag-Leffler (1846-1927)
- Implicit function theorem
- Analytic extension of functions (analytic continuation)

#### Big picture idea:

solutions "continue" locally under small parameter changes





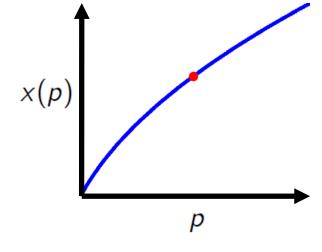


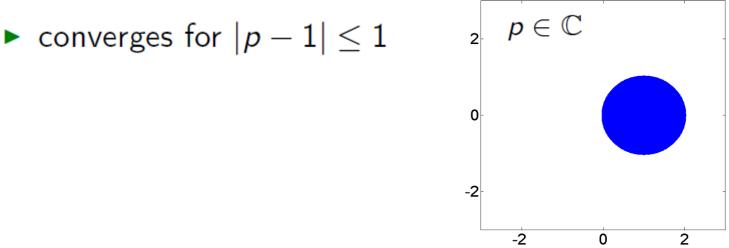
# Continuation

$$f(x; p) = x^2 - p = 0$$

Locally near p=1:

$$x(p) = \sqrt{p} = \sum_{n=0}^{\infty} \frac{(-1)^n (2n)!}{4^n (1-2n)(n!)^2} (p-1)^n$$









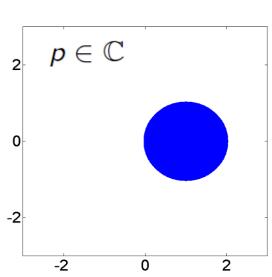
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Locally near p = 1:

$$x(p) = \sqrt{p} = \sum_{n=0}^{\infty} \frac{(-1)^n (2n)!}{4^n (1-2n)(n!)^2} (p-1)^n$$

▶ converges for  $|p-1| \le 1$ 



Use continuation to extend beyond this domain.

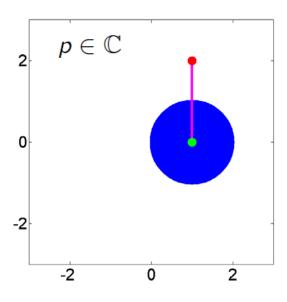




# Continuation

$$f(x; p) = x^2 - p = 0$$

Continue the solution x = 1 at p = 1 to p = 1 + 2i.





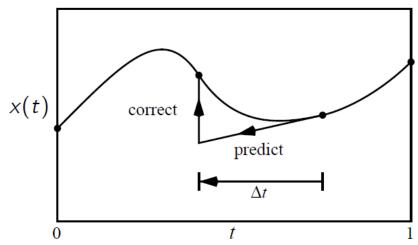


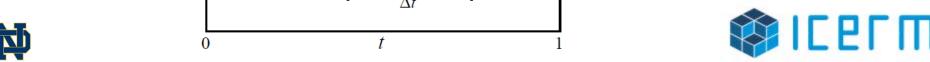
Numerically track along the path x(t) satisfying H(x(t), t) = 0:

▶ (Predictor) Estimate  $x(t + \Delta t)$  from x(t) by discretizing using the Davidenko differential equation (1953):

$$H=0 \longrightarrow \frac{d}{dt}H=0 \longrightarrow \dot{x}(t)=-J_xH(x(t),t)^{-1}J_tH(x(t),t)$$

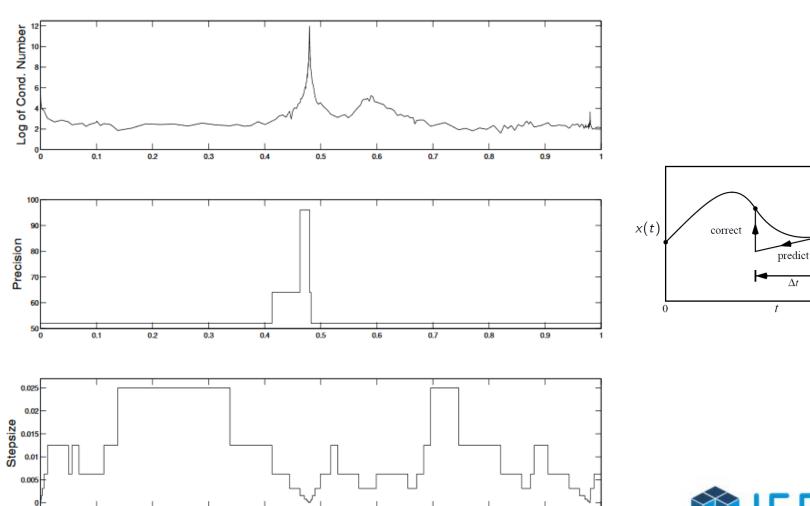
- Constant, Euler, Heun, Runge-Kutta, Runge-Kutta-Fehlberg, ....
- ▶ (Corrector) for each t, apply Newton's method to  $H(\bullet, t) = 0$





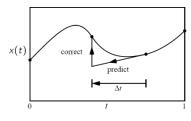
Locally adapt both stepsize and floating-point precision:

► Bates-H.-Sommese-Wampler (2008,2009), Bates-H.-Sommese (2011)









Certified tracking (select stepsize to guarantee to track path):

Shub-Smale ("Bézout series" 1990s), Beltran-Leykin (2011,2012),
 H.-Liddell (2016), Xu-Burr-Yap (2018), ...

Smale's 17<sup>th</sup> problem: polynomial time to compute a root

▶ Beltran-Pardo (2009, 2011), Cucker-Bürgisser (2011), Lairez (2017)



Pierre Lairez: 2017 SIAG/AG Early Career Prize



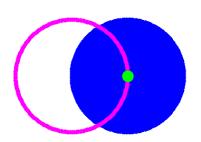


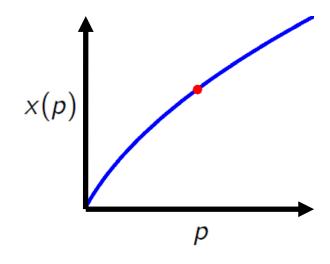
# Continuation

$$f(x; p) = x^2 - p = 0$$

Track around a loop:  $x(e^{i\theta})$ 

$$p \in \mathbb{C}$$





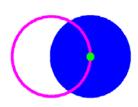




 $p \in \mathbb{C}$ 

$$f(x; p) = x^2 - p = 0$$

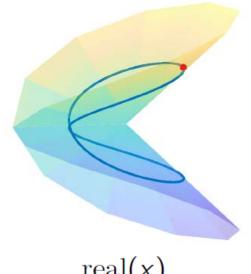
Track around a loop:  $x(e^{i\theta})$ 



• 
$$\theta = 0$$
:  $x = 1$ 

▶ 
$$\theta = 2\pi$$
:  $x = -1$ 

• 
$$\theta = 4\pi$$
:  $x = 1$ 







imag(x)

cycle number = winding number = 2

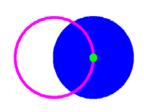




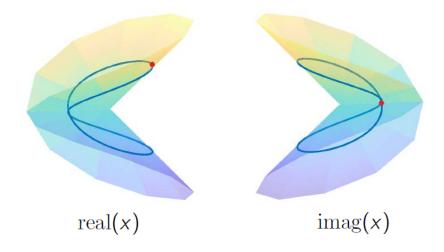
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Track around a loop:  $x(e^{i\theta})$ 

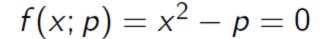


- monodromy action: permutation of solutions along loop
  - compute other solutions
    - Duff-Hill-Jensen-Lee-Leykin-Sommars (2018),
       Bliss-Duff-Leykin-Sommars (2018)
  - decompose solution sets

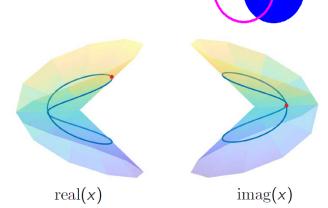












- Cauchy integral theorem: computing singular endpoints
  - cycle number c
  - sufficiently small radius r > 0

$$x(0) = \frac{1}{2\pi c} \int_0^{2\pi c} x(re^{i\theta}) d\theta$$

Cauchy endgame: Morgan-Sommese-Wampler (1991)





Find all isolated solutions of

$$f(x) = \begin{vmatrix} f_1(x_1, \dots, x_n) \\ \vdots \\ f_n(x_1, \dots, x_n) \end{vmatrix} = 0$$











$$f(x) = \begin{vmatrix} f_1(x_1, \dots, x_n) \\ \vdots \\ f_n(x_1, \dots, x_n) \end{vmatrix} = 0$$

Homotopy continuation requires (Morgan-Sommese (1989)):

- 1. parameters to "continue"
  - think of f as a member of a family  $\mathcal{F}$

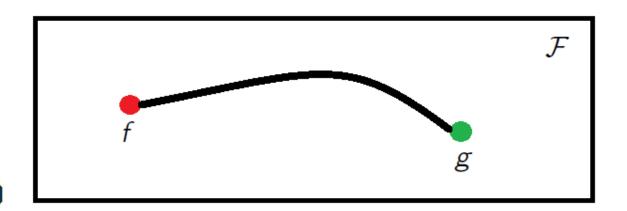






Homotopy continuation requires (Morgan-Sommese (1989)):

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- 2. homotopy that describes the deformation of the parameters
  - $\triangleright$  construct a deformation inside of  $\mathcal{F}$  that ends at f

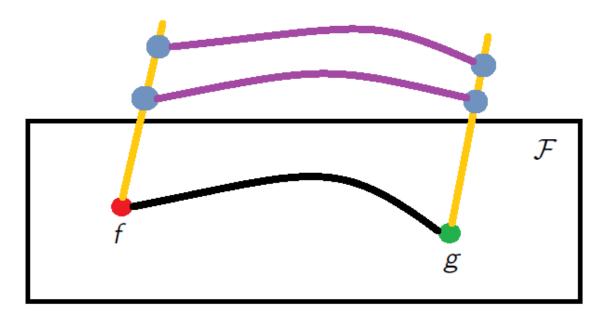






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- 1. parameters to "continue"
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- 2. homotopy that describes the deformation of the parameters
  - $\triangleright$  construct a deformation inside of  $\mathcal{F}$  that ends at f
- 3. start points to track along paths as parameters deform
  - parallelize computation track each path independently





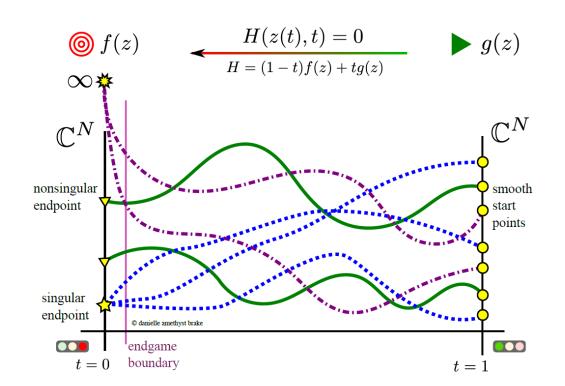


## Theorem

# Isolated Solutions

For properly constructed homotopies, with finite endpoints  $S \subset \mathbb{C}^n$ :

- each isolated solution is contained in S
  - ▶ in fact, S contains a point on every connected component
- for square systems, multiplicity = number of paths if isolated.

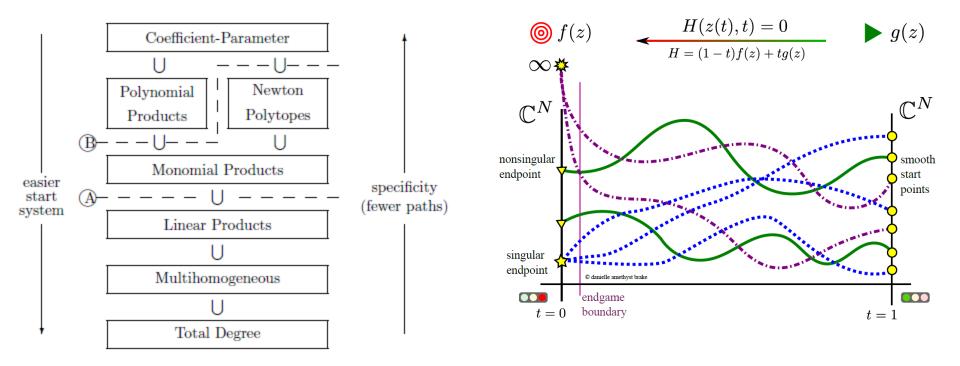






Art in the construction of family  $\mathcal{F}$ :

- number of start points
- ease to compute start points



Each method is sharp for generic members of  $\mathcal{F}$ .





# **Isolated Solutions**

$$f = \begin{bmatrix} x^2 + 2x - 8 \\ xy + 2x + 4y - 3 \end{bmatrix}$$





# **Isolated Solutions**

$$f = \begin{bmatrix} x^2 + 2x - 8 \\ xy + 2x + 4y - 3 \end{bmatrix}$$

Bézout family (total degree):

$$\mathcal{F} = \left\{ \begin{bmatrix} g_1(x,y) \\ g_2(x,y) \end{bmatrix} : \deg g_i = 2 \right\} \qquad g = \begin{bmatrix} x^2 - 1 \\ y^2 - 1 \end{bmatrix}$$

Number of paths = number of isolated solutions for g: 4

$$H = (1 - t) \cdot f + \gamma t \cdot g$$

- $ightharpoonup \gamma \in \mathbb{C}$  is used to create a general deformation
  - avoid singularities that arise from tracking over real numbers





$$f = \left[ \begin{array}{c} x^2 + 2x - 8 \\ xy + 2x + 4y - 3 \end{array} \right]$$

Bézout family (total degree):

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Number of paths = number of isolated solutions for g: 4

Bertini

input

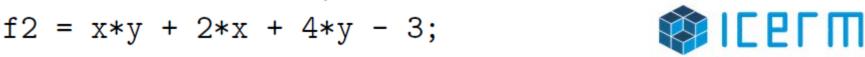
finite\_solutions

2.000000000000000e+00 0.00000000000000e+00

-1.666666666666667e-01 0.000000000000000e+00

variable\_group x,y; function f1,f2;

 $f1 = x^2 + 2*x - 8;$ 



$$f = \left[ \begin{array}{c} x^2 + 2x - 8 \\ xy + 2x + 4y - 3 \end{array} \right]$$

Multihomogeneous Bézout family (Morgan-Sommese (1987)):

$$\mathcal{F} = \left\{ \left[ egin{array}{l} g_1(x) \ g_2(x,y) \end{array} 
ight] : egin{array}{l} \deg_x g_1 = 2, \ \deg_x g_2 = \deg_y g_2 = 1 \end{array} 
ight\}$$

$$g = \begin{bmatrix} x^2 - 1 \\ (x - 2)(y - 1) \end{bmatrix}$$
  $H = (1 - t) \cdot f + \gamma t \cdot g$ 

Number of paths = number of isolated solutions for g: 2





$$f = \left[ \begin{array}{c} x^2 + 2x - 8 \\ xy + 2x + 4y - 3 \end{array} \right]$$

► Multihomogeneous Bézout family (Morgan-Sommese (1987)):

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ight\}$$

Number of paths = number of isolated solutions for g: 2

Bertini

input variable\_group x;
 variable\_group y;

function f1,f2;

$$f1 = x^2 + 2*x - 8;$$
  
 $f2 = x*y + 2*x + 4*y - 3;$ 





$$f = \begin{bmatrix} x^2 + 2x - 8 \\ xy + 2x + 4y - 3 \end{bmatrix}$$

Polyhedral (BKK, Huber-Sturmfels (1995)):

$$\mathcal{F} = \left\{ \begin{bmatrix} a_1 x^2 + a_2 x + a_3 \\ a_4 x y + a_5 x + a_6 y + a_7 \end{bmatrix} : a_i \in \mathbb{C} \right\}$$

$$g = \begin{bmatrix} x^2 - 1 \\ y - 1 \end{bmatrix} \qquad H = (1 - t) \cdot f + \gamma t \cdot g$$

Number of paths = number of isolated solutions for g: 2





# **Isolated Solutions**

$$f = \begin{bmatrix} x^2 + 2x - 8 \\ xy + 2x + 4y - 3 \end{bmatrix}$$

Extra structure in the coefficients of f.

$$\mathcal{F} = \left\{ p(x, y; a) = \begin{bmatrix} x^2 - (a_1 + a_2)x + a_1 a_2 \\ (x - a_1)y + a_3 x + a_4 \end{bmatrix} : a_i \in \mathbb{C} \right\}$$

$$g = \left[ \begin{array}{c} x^2 - 1 \\ (x - 1)y - 1 \end{array} \right]$$

Number of paths = number of isolated solutions for g: 1





$$f = \left[ \begin{array}{c} x^2 + 2x - 8 \\ xy + 2x + 4y - 3 \end{array} \right]$$

$$\mathcal{F} = \left\{ p(x, y; a) = \left[ \begin{array}{c} x^2 - (a_1 + a_2)x + a_1 a_2 \\ (x - a_1)y + a_3 x + a_4 \end{array} \right] : a_i \in \mathbb{C} \right\}$$

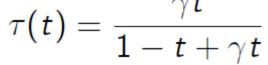
$$g = \left[ \begin{array}{c} x^2 - 1 \\ (x - 1)y - 1 \end{array} \right]$$

Since  $\mathcal{F}$  is no longer linear, use a parameter homotopy:

$$H = p(x, y; a(t))$$

where  $a(t) = (1 - \tau(t))(-4, 2, 2, -3) + \tau(t)(1, -1, 0, -1)$ 



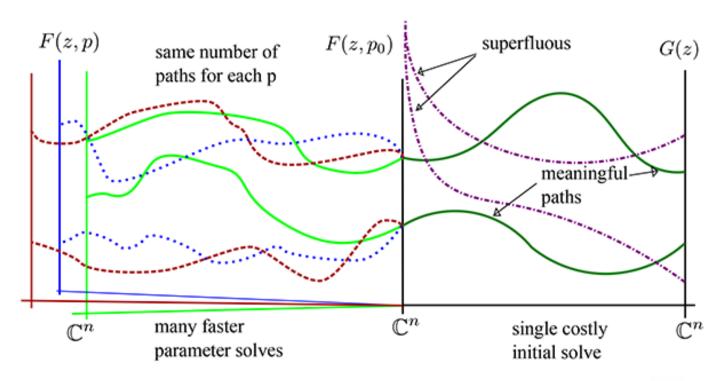




# Isolated Solutions

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#### Some software options:

- ▶ Bertini
- ▶ Bertini.m2
- ► Hom4PS
- HomotopyContinuation.jl
- MonodromySolver
- ► NAG4M2
- Paramotopy
- PHCpack

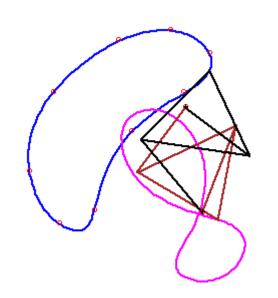
Visitors to ICERM: Bates, Brake, Chen, Duff, Hill, Lee, Leykin, Rodriguez, Sommars, Sommese, Wampler, ...

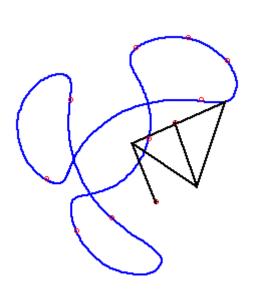


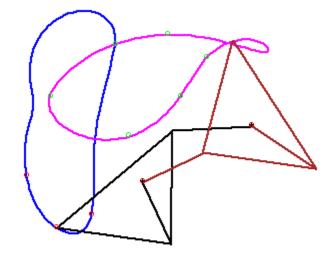


Example (Alt's problem (1923))

Find all 4-bar linkages whose coupler curve passes through 9 given general points in the plane.

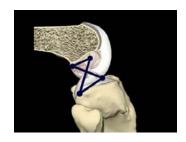
















Example (Alt's problem (1923))

Find all 4-bar linkages whose coupler curve passes through 9 given general points in the plane.

▶  $8652 = 6 \cdot 1442$  (Wampler-Morgan-Sommese (1992))

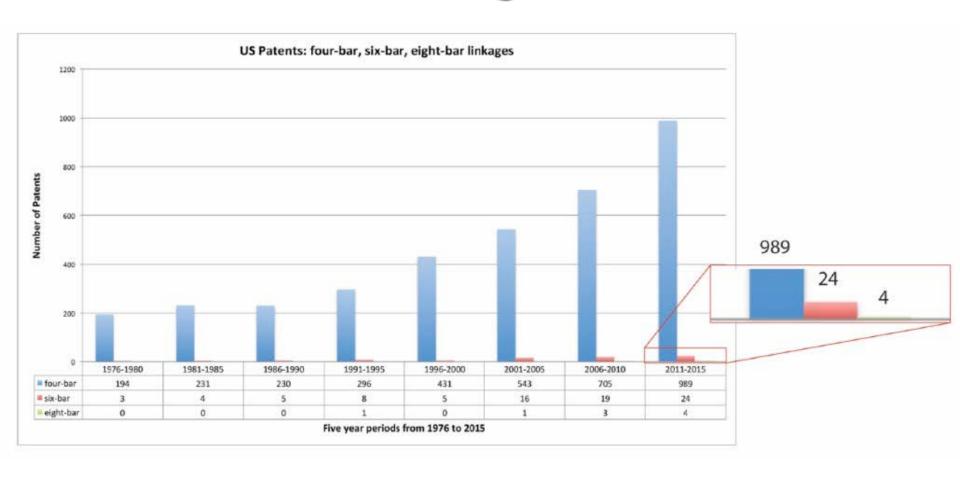
Their polynomial system: 4 quadratics and 8 quartics

Bézout	1,048,576	$= 2^4 \cdot 4^8$
M-hom Bézout	286,720	$=2^{12}\cdot \binom{8}{4}$
Polyhedral	79,135	
Product decomp.	18,700	
Actual	8,652	





# Mechanical Design 101 MECHANICAL DESIGN EDUCATIONAL RESOURCE

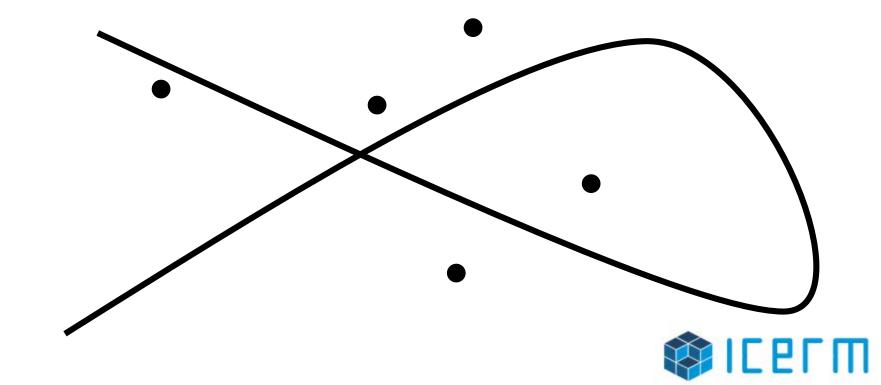






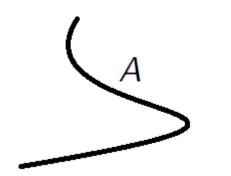
Describe all solutions of

$$f(x) = \begin{bmatrix} f_1(x_1, \dots, x_n) \\ \vdots \\ f_k(x_1, \dots, x_n) \end{bmatrix} = 0$$





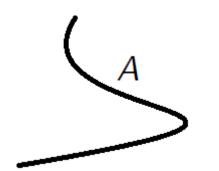
How to represent an irreducible algebraic variety A on a computer?







How to represent an irreducible algebraic variety A on a computer?



- ▶ algebraic: prime ideal  $I(A) = \{g \mid g(a) = 0 \text{ for all } a \in A\}$ 
  - ▶ Hilbert Basis Theorem (1890): there exists  $f_1, \ldots, f_k$  such that

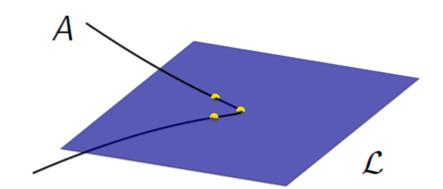
$$I(A) = \langle f_1, \ldots, f_k \rangle$$





How to represent an irreducible algebraic variety A on a computer?

- ▶ geometric: witness set  $\{f, \mathcal{L}, W\}$  where
  - f is polynomial system where A is an irred. component of  $\mathcal{V}(f)$
  - $\mathcal{L}$  is a linear space with  $\operatorname{codim} \mathcal{L} = \dim A$
  - ▶  $W = \mathcal{L} \cap A$  where  $\#W = \deg A$



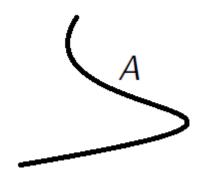
- Witness sets "localize" computations to A effectively ignoring other irreducible components
- ightharpoonup Sample points from A by moving the linear slice  $\mathcal{L}$





## Witness Set

$$A = \{[s^3, s^2t, st^2, t^3] \mid [s, t] \in \mathbb{P}^1\} \subset \mathbb{P}^3$$
 – twisted cubic curve



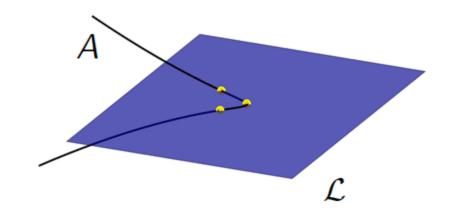




## Witness Set

$$A = \{[s^3, s^2t, st^2, t^3] \mid [s, t] \in \mathbb{P}^1\} \subset \mathbb{P}^3$$
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- ▶  $\{f, \mathcal{L}, W\}$  where



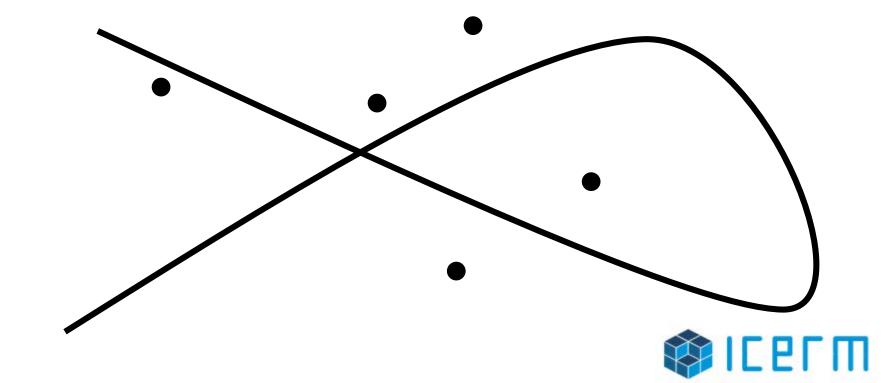
- ▶  $\mathcal{L} = \{ [x_0, x_1, x_2, x_3] \in \mathbb{P}^3 \mid 6x_0 6x_1 2x_2 + x_3 = 0 \} \subset \mathbb{P}^3$ ▶  $\operatorname{codim} \mathcal{L} = \dim A = 1$
- $W = \left\{ \begin{array}{l} [1, 3.2731, 10.7130, 35.0644], \\ [1, 0.8596, 0.7389, 0.6351], \\ [1, -2.1326, 4.5481, -9.6995] \end{array} \right\}$ 
  - ▶ deg A = 3





Numerical irreducible decomposition:

compute a witness set for each irreducible component





$$f = \left[ \begin{array}{c} x_1^2 - x_0 x_2 \\ x_1 x_2 - x_0 x_3 \end{array} \right]$$

Witness Set

```
Bertini
                input
CONFIG
```

```
TrackType: 1;
```

END;

```
INPUT
```

```
hom_variable_group x0,x1,x2,x3;
function f1,f2;
f1 = x1^2 - x0*x2;
f2 = x1*x2 - x0*x3;
END;
```

Dimension 1: 2 classified components

degree 1: 1 component degree 3: 1 component





Reduce to codimension = # equations via randomization:

## Theorem (Bertini)

Let  $f: \mathbb{C}^n \to \mathbb{C}^N$  and  $A \subset \mathcal{V}(f) \subset \mathbb{C}^n$  be an irreducible component with codim A = c. If  $R \in \mathbb{C}^{c \times N}$  is general, then

- ightharpoonup A is an irreducible component of  $V(R \cdot f)$
- $\triangleright$   $V(R \cdot f) \setminus V(f)$  is either empty or smooth of codimension c.





Reduce to codimension = # equations via randomization:

## Theorem (Bertini)

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- ▶  $V(R \cdot f) \setminus V(f)$  is either empty or smooth of codimension c.

### Example

For general 
$$R \in \mathbb{C}^{2 \times 3}$$
 and  $f = \begin{bmatrix} x_1^2 - x_0 x_2 \\ x_1 x_2 - x_0 x_3 \\ x_2^2 - x_1 x_3 \end{bmatrix}$ ,

 $V(R \cdot f) = \text{twisted cubic} + \text{line}$ 





## Witness Set

$$f = \begin{bmatrix} (x - y)(\hat{x} - \hat{y}) \\ (x - y)(\hat{a}\hat{x} - 2\hat{a}\hat{y} + 2\hat{b}\hat{x} - \hat{b}\hat{y}) \\ (\hat{x} - \hat{y})(ax - 2ay + 2bx - by) \\ \hat{a}\hat{b}(x - y)(\hat{a}\hat{y} - \hat{b}\hat{x}) \\ ab(\hat{x} - \hat{y})(ay - bx) \\ \vdots \\ (a\hat{b}\hat{x}y - \hat{a}bx\hat{y})(a\hat{b}\hat{x} - \hat{a}\hat{b}x - \hat{a}b\hat{y} + \hat{a}\hat{b}y - a\hat{x}\hat{y} + \hat{a}x\hat{y} + b\hat{x}\hat{y} - \hat{b}\hat{x}y) \\ (a\hat{b}\hat{x}y - \hat{a}bx\hat{y})(ab\hat{x} - \hat{a}bx - ab\hat{y} + a\hat{b}y - a\hat{x}y + \hat{a}xy + bx\hat{y} - \hat{b}xy) \end{bmatrix}$$

15 polynomials in 8 variables  $a, b, x, y, \hat{a}, \hat{b}, \hat{x}, \hat{y}$ 

For general  $R \in \mathbb{C}^{8 \times 15}$ :

- $\blacktriangleright V(R \cdot f) \setminus V(f)$  consists of finitely many points
  - ▶ all nonsingular with respect to  $R \cdot f = 0$





## Witness Set

$$f = \begin{bmatrix} (x - y)(\hat{x} - \hat{y}) \\ (x - y)(\hat{a}\hat{x} - 2\hat{a}\hat{y} + 2\hat{b}\hat{x} - \hat{b}\hat{y}) \\ (\hat{x} - \hat{y})(ax - 2ay + 2bx - by) \\ & \hat{a}\hat{b}(x - y)(\hat{a}\hat{y} - \hat{b}\hat{x}) \\ & ab(\hat{x} - \hat{y})(ay - bx) \\ & \vdots \\ & (a\hat{b}\hat{x}y - \hat{a}bx\hat{y})(a\hat{b}\hat{x} - \hat{a}\hat{b}x - \hat{a}b\hat{y} + \hat{a}\hat{b}y - a\hat{x}\hat{y} + \hat{a}x\hat{y} + b\hat{x}\hat{y} - \hat{b}\hat{x}y) \\ & (a\hat{b}\hat{x}y - \hat{a}bx\hat{y})(ab\hat{x} - \hat{a}bx - ab\hat{y} + a\hat{b}y - a\hat{x}y + \hat{a}xy + bx\hat{y} - \hat{b}xy) \end{bmatrix}$$

15 polynomials in 8 variables  $a, b, x, y, \hat{a}, \hat{b}, \hat{x}, \hat{y}$ 

For general  $R \in \mathbb{C}^{8 \times 15}$ :

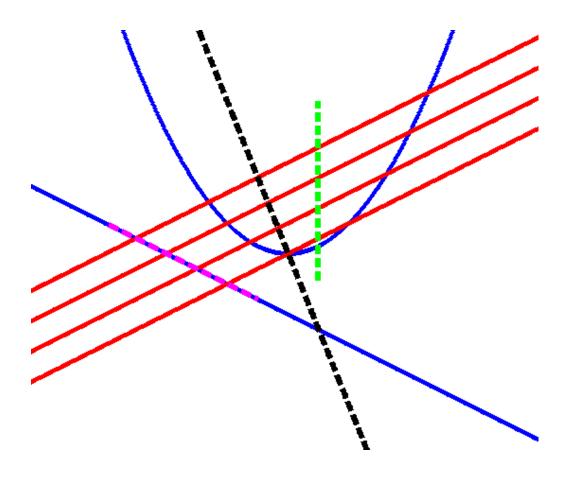
- ▶  $V(R \cdot f) \setminus V(f)$  consists of finitely many points
  - ▶ all nonsingular with respect to  $R \cdot f = 0$
- ▶ Using Bertini:  $|V(R \cdot f) \setminus V(f)| = 8652$ 
  - Proving this would complete proof of Alt's problem





Given  $W \subset V(f) \cap \mathcal{L}$ , how to test that  $W = \mathcal{L} \cap A$  for some variety  $A \subset V(f)$ ?

Trace test: centroid moves linearly as slices moves in parallel

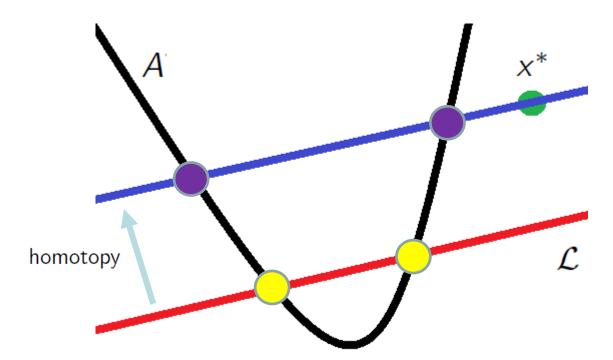






Many other numerical algebraic geometric computations can be performed starting from witness sets, such as:

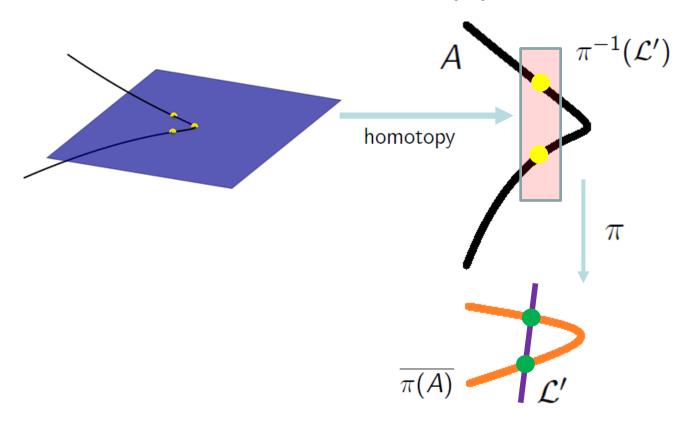
- ▶ membership testing: is  $x^* \in A$ ?
  - ▶ decide if  $g(x^*) = 0$  for every  $g \in I(A)$  without knowing I(A)







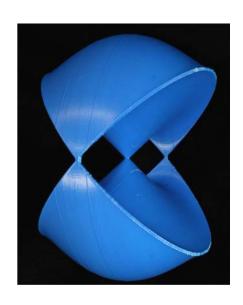
- projection:  $\overline{\pi(A)}$ 
  - perform computations on  $\overline{\pi(A)}$  without knowing any polynomials that vanish on  $\overline{\pi(A)}$

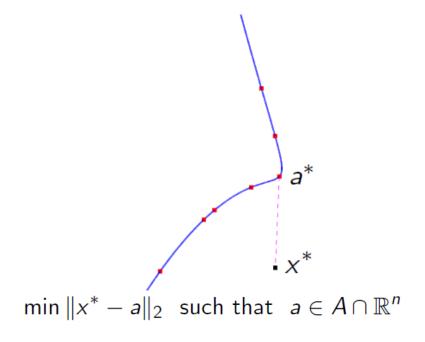






- intersection: A ∩ B
  - special case is regeneration
    - $\mathcal{V}(f_1,\ldots,f_k,f_{k+1})=\mathcal{V}(f_1,\ldots,f_k)\cap\mathcal{V}(f_{k+1})$  via witness sets
  - ightharpoonup compute  $A_{\rm sing}$
  - compute critical points of optimization problem



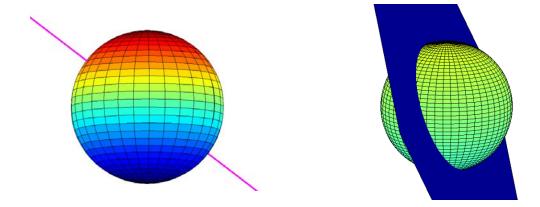






Test other algebraic properties of A

- ▶ is A arithmetically Cohen Macaulay?
- ▶ is A arithmetically Gorenstein?
- ▶ is A a complete intersection?







# Summary

Numerical algebraic geometry provides a toolbox for solving polynomial systems.

- "If a problem was easy, someone else would have solved it."
  - Gröbner basis computation probably did not terminate
- think carefully about what information you want/need
- art in building efficient homotopies that incorporate structure
- preconditioning is important
  - transform problem into form suitable for num. computations





# Thank You!



